15-213 Recitation 13

Introduction to Computer Systems

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Proxy Lab

- Due 4th December, 2013
- No late days for the assignment!
- We will have interviews with each of you
 - NOT a group based assignment
 - Interviews are on the Reading Day, the schedule will be out soon!

Proxy lab

- Three main steps
 - Step 1: Implement a sequential proxy
 - Step 2: Make the proxy concurrent
 - Step 3: Implement a web-cache

Step 1: Implement a sequential proxy

- Check out last recitation!
- Hopefully you are done with this, or at least started and are close

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- So you can share data very easily between threads
- With great power, comes great responsibility (and headaches)
 - Possibility of race conditions

```
#include "csapp.h"
```

static volatile int global = 0;

```
void *thread(void *vargp) {
    int i;
    for (i = 0; i < 100; i++) {
      global++;
    }
    return NULL;
}</pre>
```

What will be the output?

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Any value from 2-200!

- Shared variable is global
- global++
 - Can we divided into three atomic operations
 1)Read the variable into a register
 2)Increment
 3)Store the variable back to memory

Normal Order of Operations $R \rightarrow I \rightarrow S$

- Lets look at some possibilities
 - 200 → This is simple, if one thread finishes completely first, and then the second thread starts
 - 100 → A little tricky, access pattern must be R R I
 S I S, where red is thread 1, and blue is thread 2.
 - 2 \rightarrow The trickiest of them all!
 - Imagine, R R I S R I S I S (Thread 2 run's 99 times)
 - I S R I S R I S I S (Thread 1 run's its remaining 99 times)

A peculiar example

• Thread's can access each other's stacks. Consider the following example:

#include<stdio.h>
#include<pthread.h>

```
void *thread(void *vargp)
{
    int myid = *(int *)vargp;
    while(1) {
         *(int *)vargp = *(int *)vargp +1;
     }
}
int main()
{
    int i;
    pthread t tid;
    pthread create(&tid, NULL, thread, (void *)&i);
    while(1)
         printf("%d\n",i);
    pthread exit(NULL);
}
```